



BONSAI Lab

 POLITECNICO DI MILANO



Survivable Virtual Topology Mapping To Provide Content Connectivity Against Double-Link Failures

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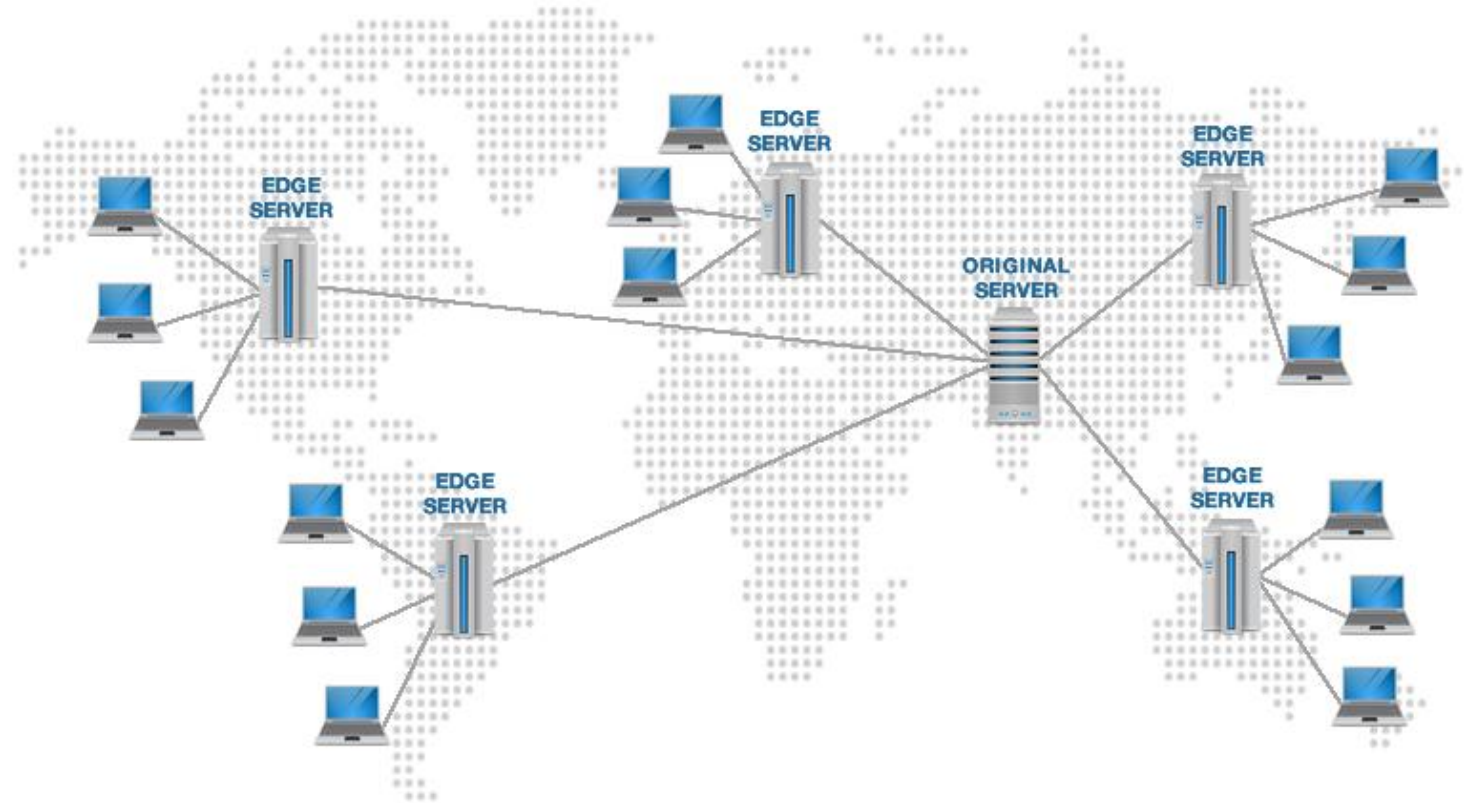


- Motivations
- Survivable virtual topology mapping problem
- Network connectivity Vs content connectivity
- Problem statement
- Case-study & results
- Conclusions & future works



CONTENT DELIVERY NETWORKS (CDNs)

- ❑ Improved user experience
- ❑ Bandwidth
 - Reduce probability of link congestion
 - Reduce energy consumption
- ❑ Reduce delay
- ❑ Availability
 - A single server is easily overloaded
- ❑ Load balancing

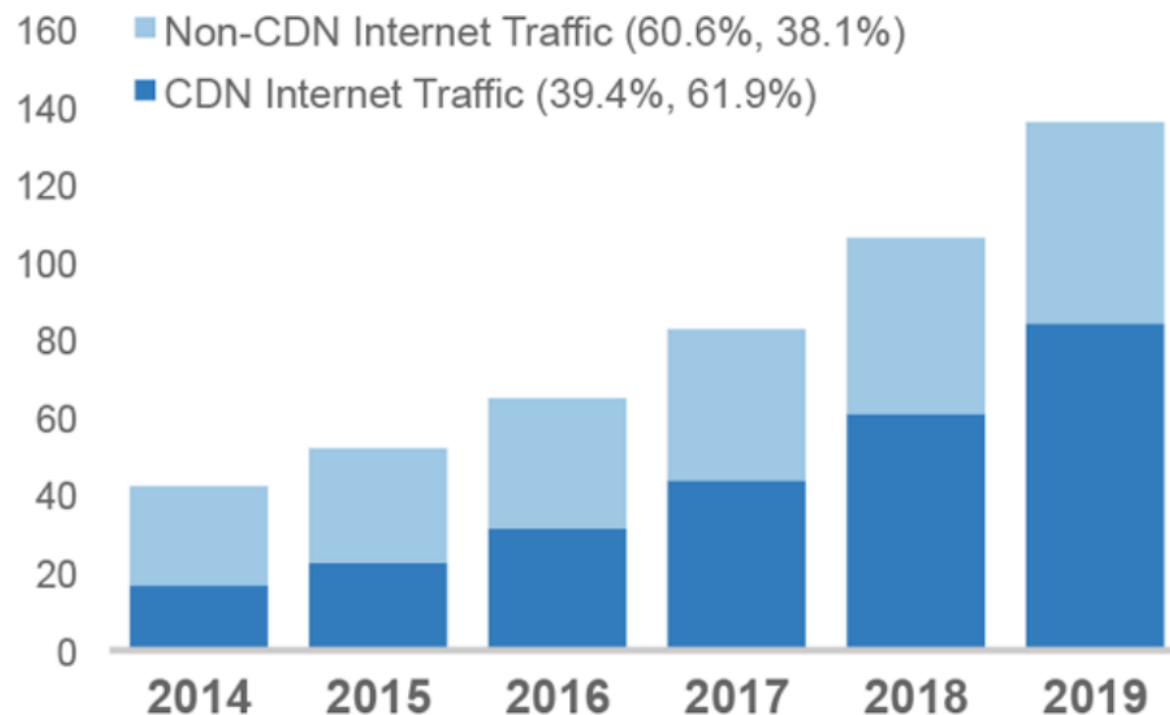




THE EVOLUTION OF CONTENT DELIVERY NETWORKS

- Social media
- Online gaming
- Electronic payment transactions
- Multimedia streaming
- ...etc

Exabytes per Month

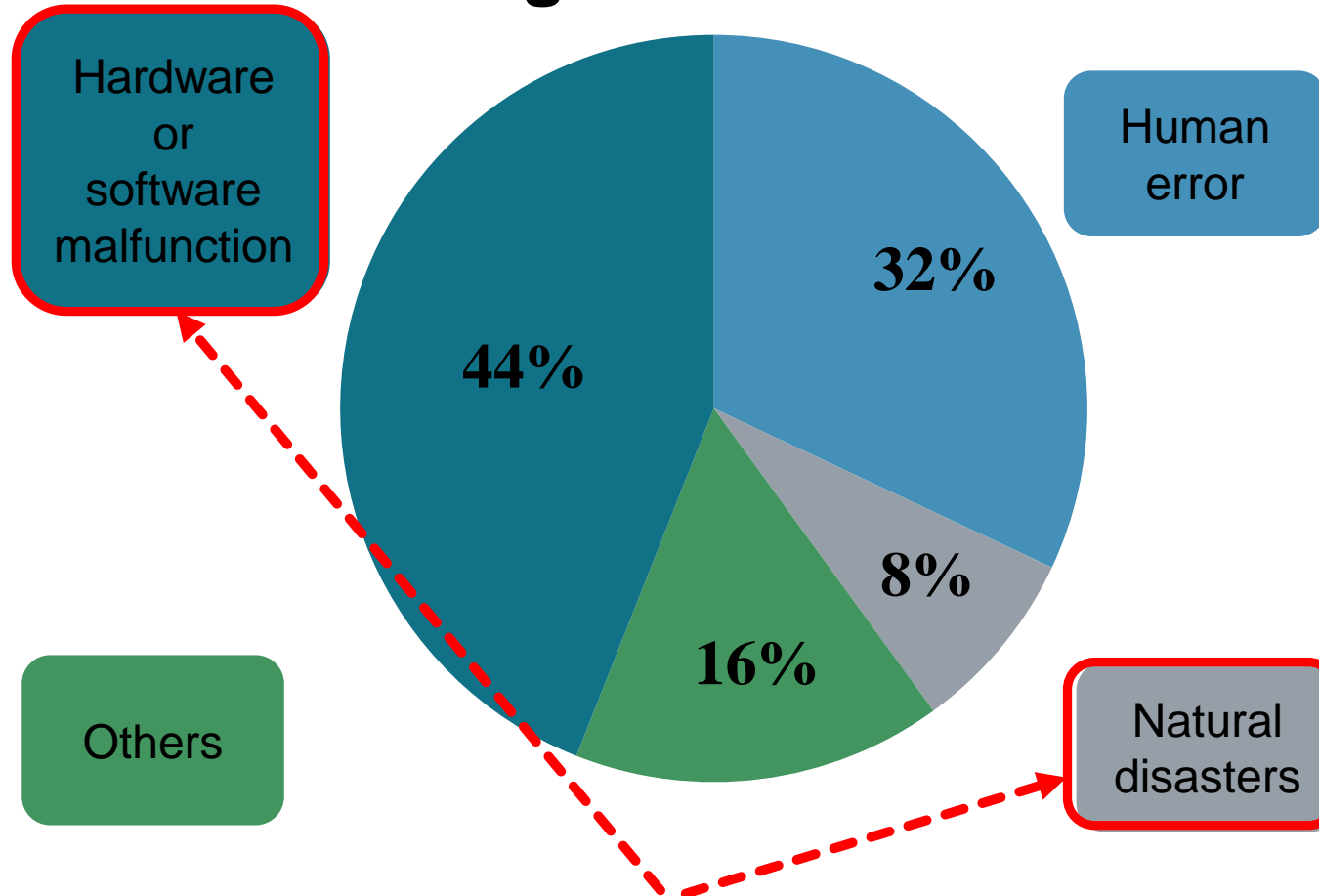


End-to-end → End-to-content

Source: Index, Cisco Visual Networking. "Forecast and Methodology, 2014–2019 White Paper, Cisco, 2015."



Leading causes of Data loss

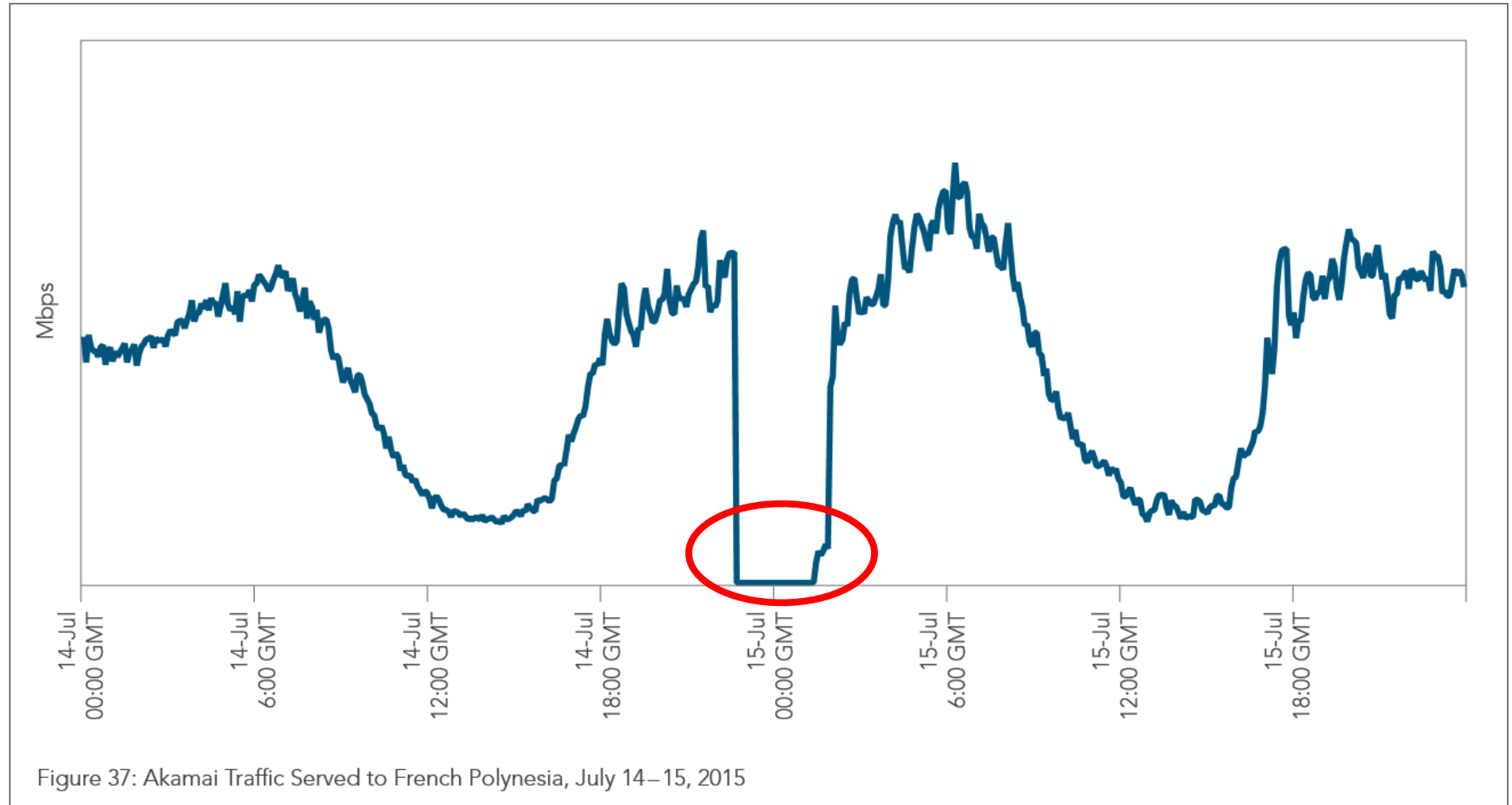


Source: <http://www.cloudwards.net/world-backup-day-2015>
http://www.paragon-software.com/survey_results.html



CHARACTERISTICS OF DISRUPTIONS FOR CDN PROVIDERS (I)

- Total loss of connectivity
- More than 6 hours to recover from failure
- Failure due to hardware malfunction

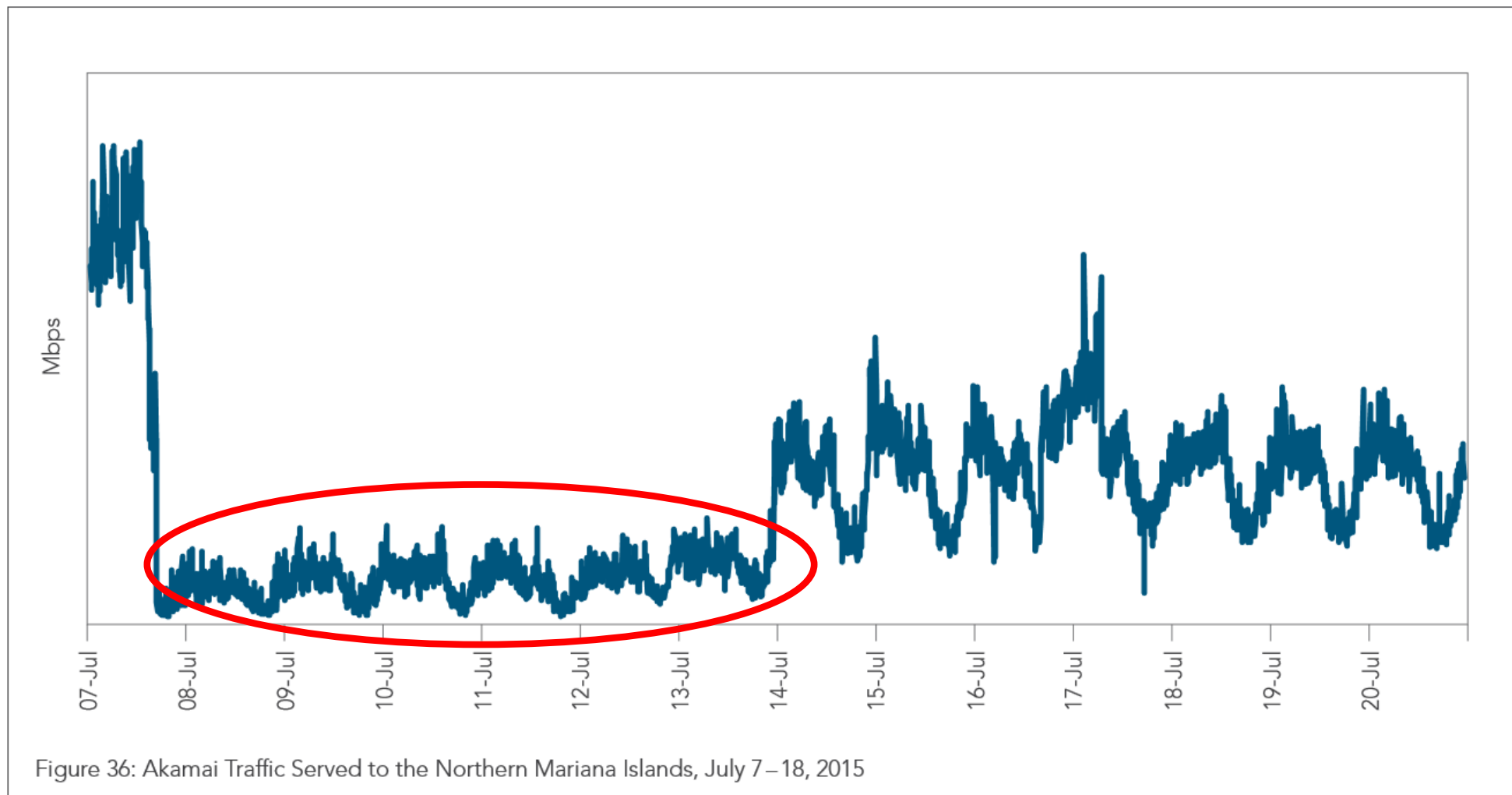


Source: Q3 2015 State of the Internet Report - AKAMAI



CHARACTERISTICS OF DISRUPTIONS FOR CDN PROVIDERS (II)

- Partial loss of connectivity
- More than 1 week to recover from failure
- Failure due fiber cut undersea

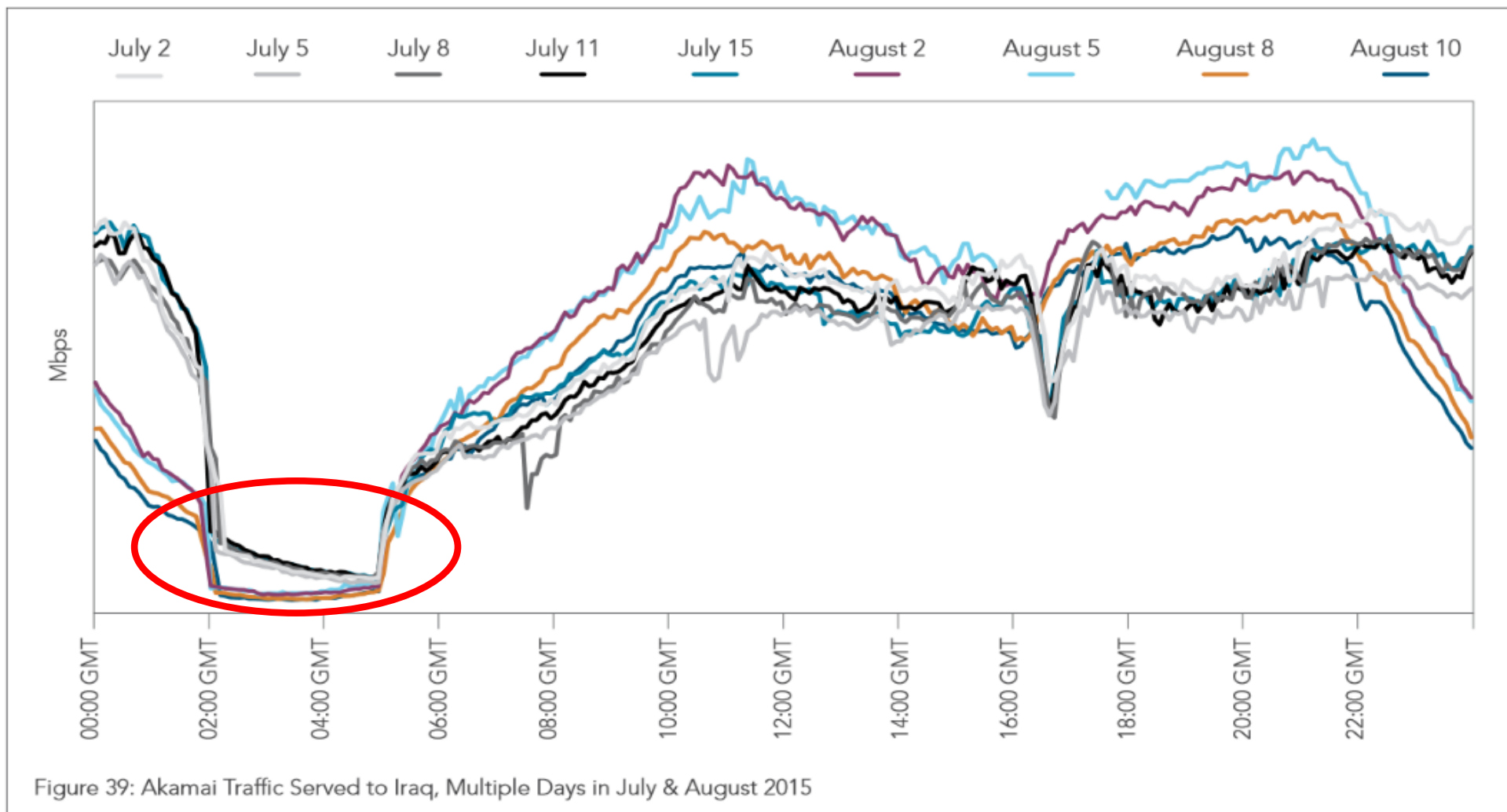


Source: Q3 2015 State of the Internet Report - AKAMAI



CHARACTERISTICS OF DISRUPTIONS FOR CDN PROVIDERS (III)

- Total loss of connectivity
- More than 3h to restore connectivity
- Government shutting down internet during national exams days



Source: Q3 2015 State of the Internet Report - AKAMAI

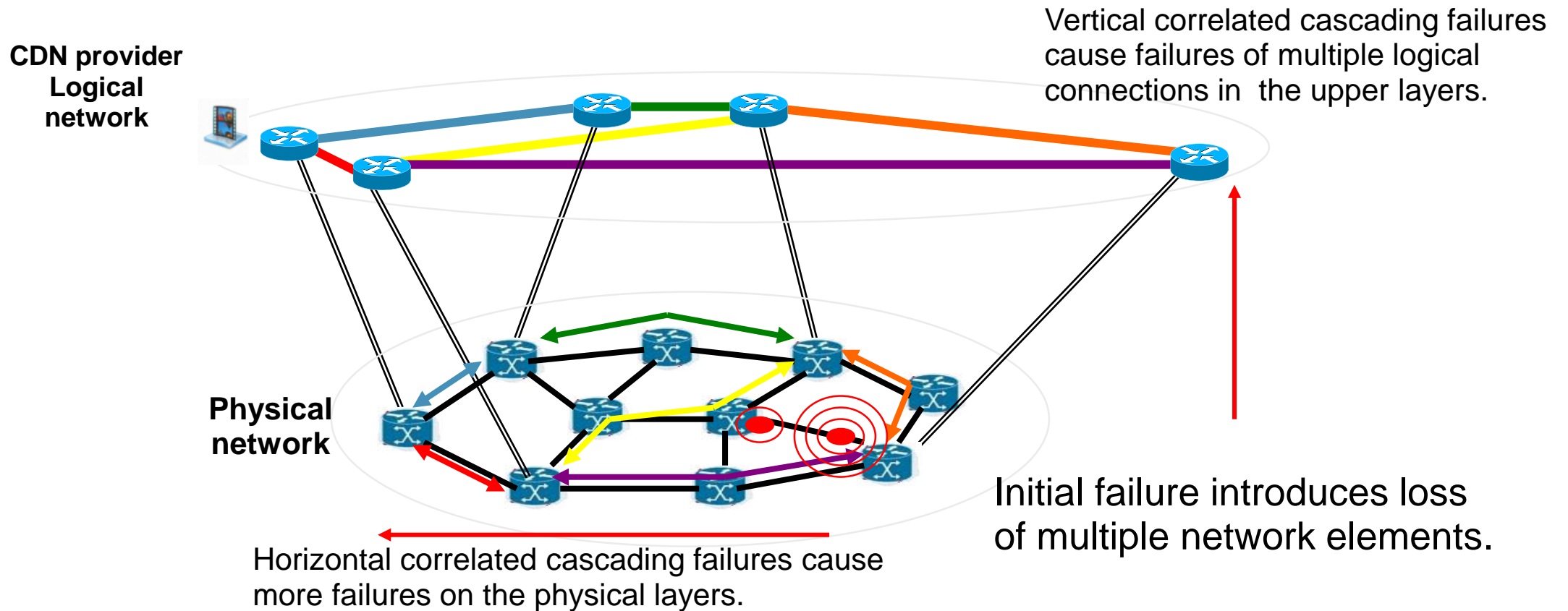


HOW CAN A CDN PROVIDER OVERCOME TOTAL LOSS OF CONNECTIVITY DURING FAILURE EVENTS?

The idea is to perform a **survivable** design combined with an **optimal content placement** to allow the CDN provider to continue the content distribution until the failure recovery from **single/multiple failures** is done.



IMPACT OF PHYSICAL LAYER FAILURES

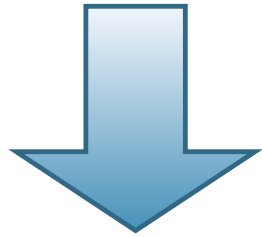




NETWORK CONNECTIVITY VS CONTENT CONNECTIVITY

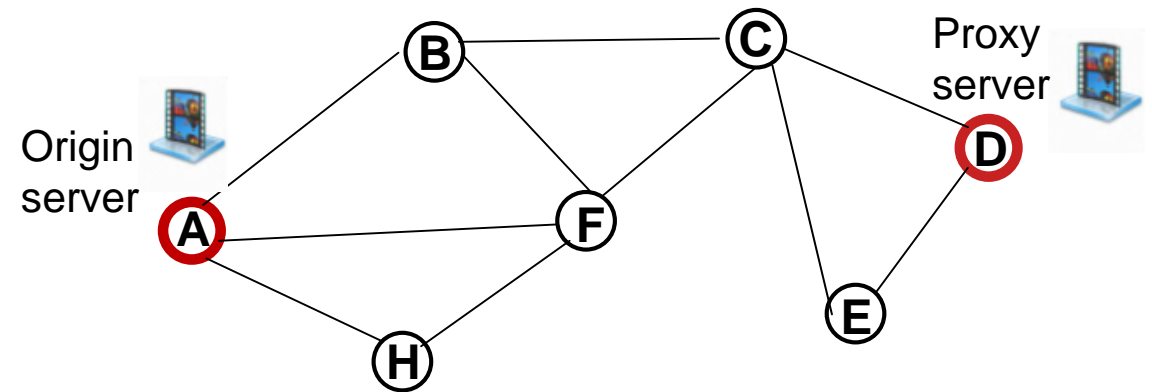
Traditional metric: **Network connectivity**

- Reachability of all nodes from any point in the network.



New metric : **Content connectivity**

- Reachability of the content from any point in the network.

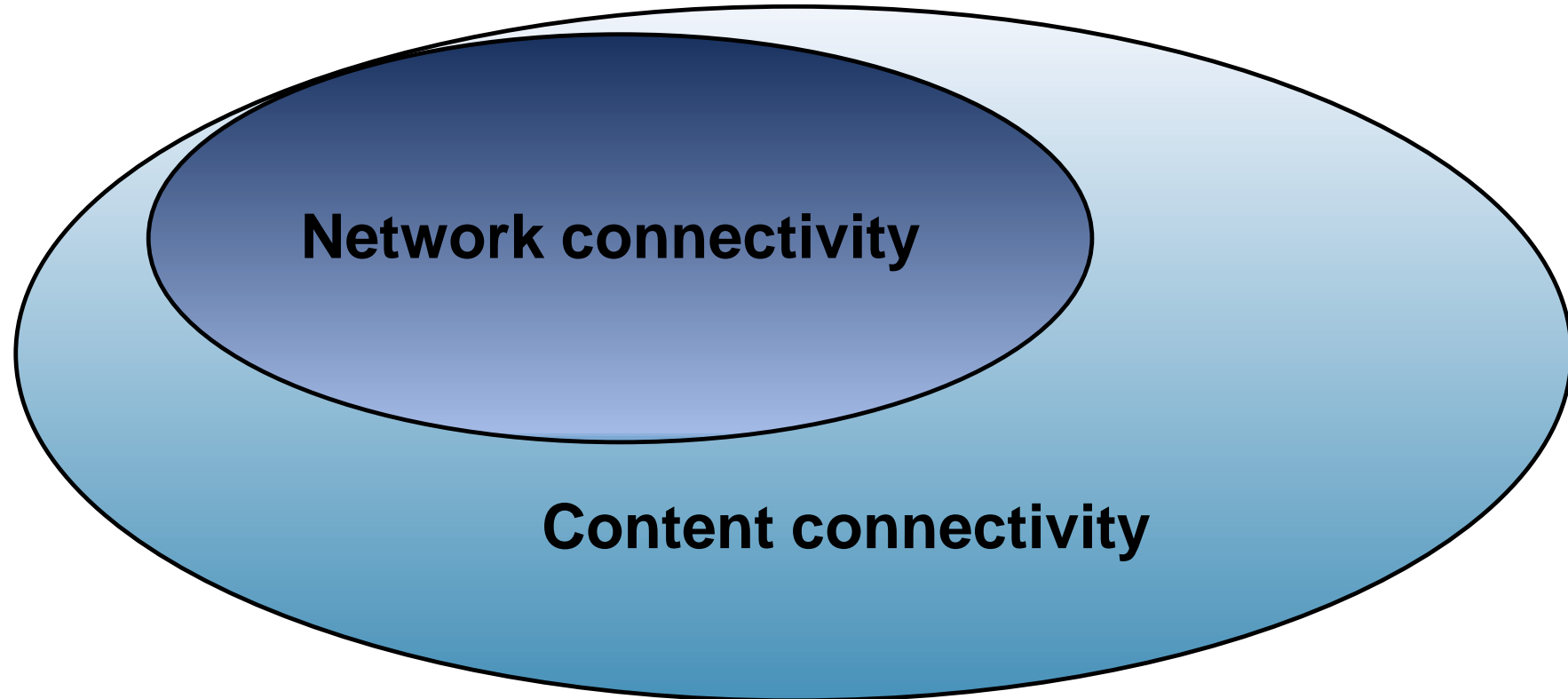


CDN provider virtual network

**“Fault-Tolerant Virtual Network Mapping to Provide Content Connectivity in Optical Networks,” M. F. Habib et al.



NETWORK CONNECTIVITY VS CONTENT CONNECTIVITY (II)



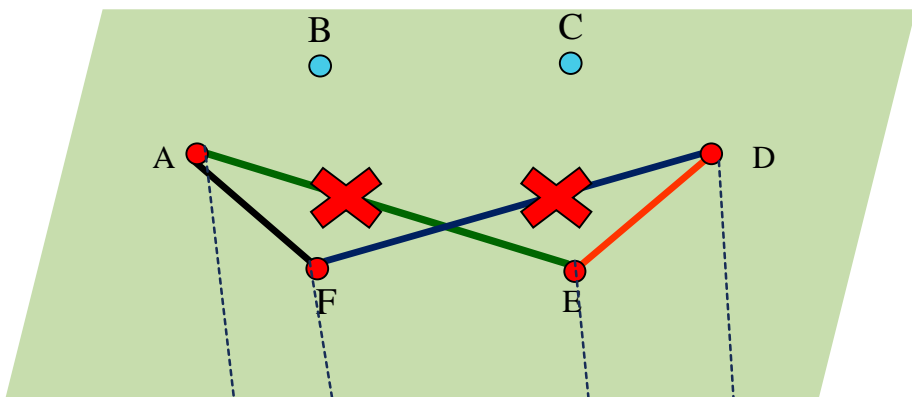


Two factors:

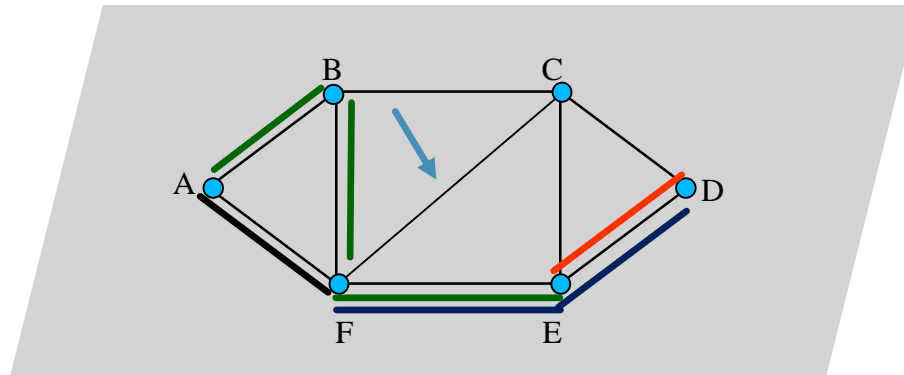
- **Survivable virtual topology mapping**
- **Optimal content placement**



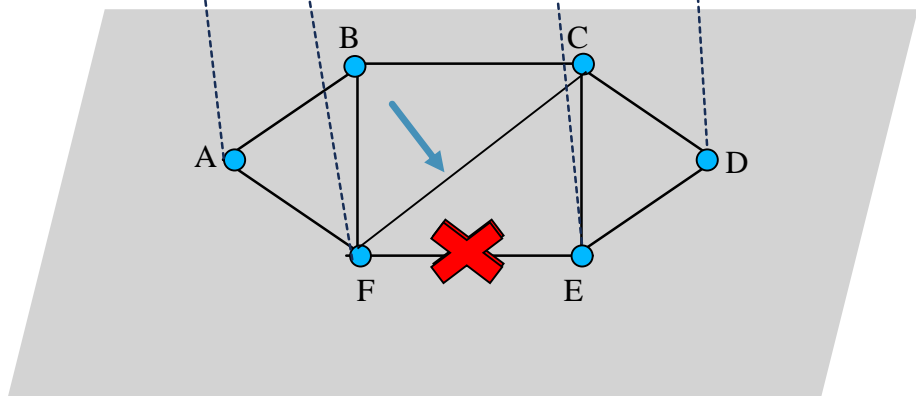
NETWORK CONNECTED SURVIVABLE MAPPING



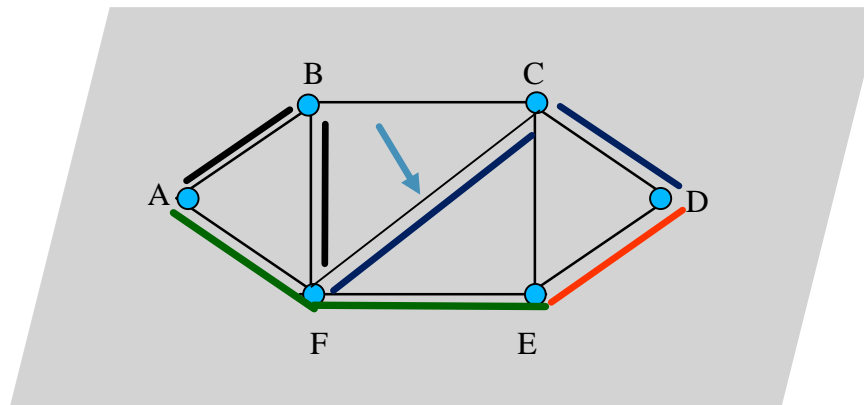
Logical Topology



Non-Survivable Mapping



Physical Topology



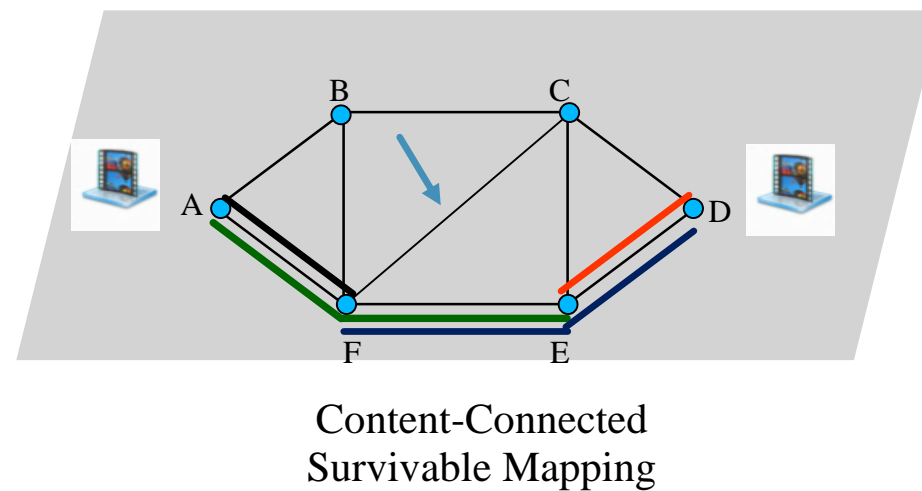
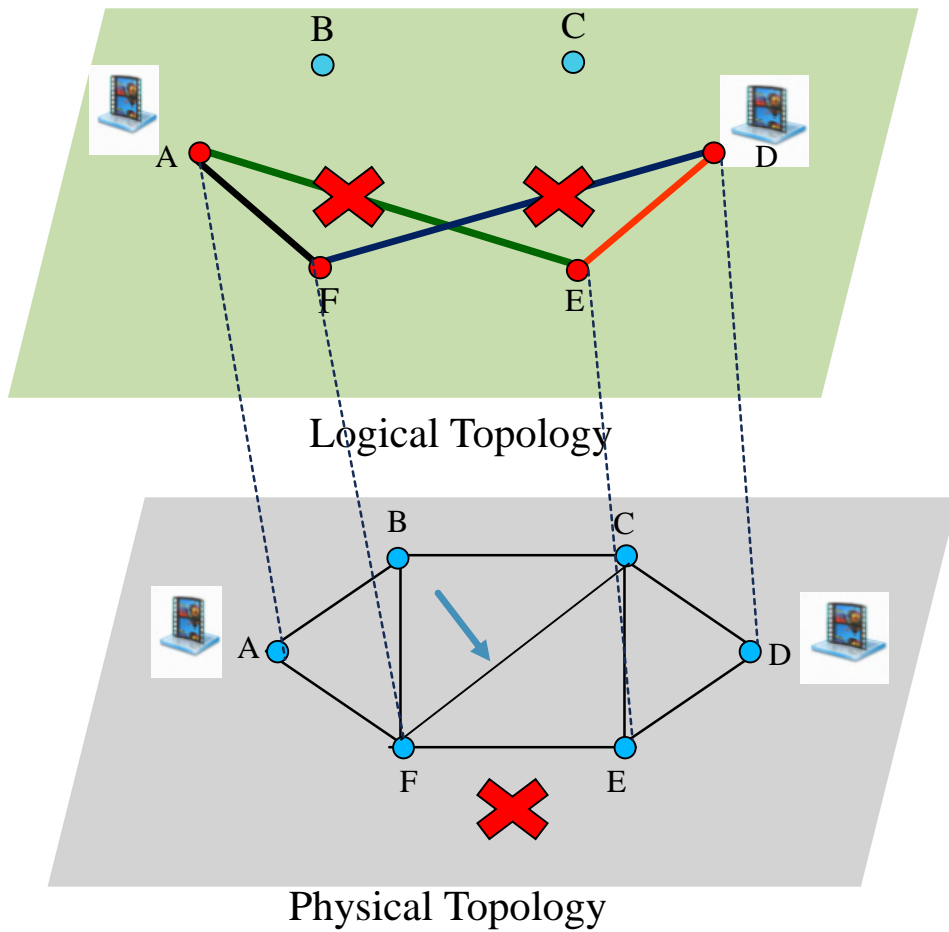
Network-Connected Survivable Mapping**

No possible survivable mapping

**Menger's theorem



CONTENT CONNECTED SURVIVABLE MAPPING



**Content
Connectivity
is still
guaranteed**



PROPOSED DESIGN SCENARIOS

Design scenario	Failure type	Survivable mapping	Optimal content placement
Content connectivity ** (CC1)	Single-link	-	✓
Network connectivity (NC1)	Single-link	✓	-
Network connectivity + Content connectivity (NC1+NC2)	Double-link	✓	✓
Content connectivity (CC2)	Double-link	-	✓
Network connectivity (NC2)	Double-link	✓	-

**“Fault-Tolerant Virtual Network Mapping to Provide Content Connectivity in Optical Networks,” M. F. Habib et al.



PROBLEM STATEMENT

□ Given

- Physical topology
- Logical topology
- Potential data center locations

□ Objective

- **Minimize the resource usage** (i.e., wavelength channels)

□ While

- Flow constraints
- Data center placement constraints
- Capacity constraints

□ Outputs

- Optimal routing
- Optimal placement of data centers

Integer Linear programming



VARIABLES

- $P_{ij}^{st} \in \{0,1\}$: $(s, t) \in E_L$ mapped over physical link $(i, j) \in E$ ← Mapping virtual links to physical links
- $A_d^{sc} \in \{0,1\}$: Virtual node $s \in V_L$ can get content $c \in C$ in data center $d \in D$ ← Content placement
- $R_d^c \in \{0,1\}$: Content $c \in C$ is replicated at data center $d \in D$
- $\alpha_z^{st} \in \{0,1\}$: Failure of physical link $z \in E$ disrupts the logical link $(s, t) \in E_L$ ← Select disturbed Virtual links
- $F_{stz}^{uc} \in \{0,1\}$: Virtual node $u \in V_L$ can get content $c \in C$ using logical link $(s, t) \in E_L$ after a failure occurs in $z \in Z$ ← Reachability to data center after single-link failures
- $T_{stz_1z_2}^{uc} \in \{0,1\}$: Virtual node $u \in V_L$ can get content $c \in C$ using logical link $(s, t) \in E_L$ after a first failure occurs in $z_1 \in E$ and a second failure occurs in $z_2 \in E$ such that $z_1 \neq z_2$ ← Reachability to data center after double-link failures



INTEGER LINEAR PROGRAMMING FORMULATION (NC1+CC2)

$$\sum_{j:(i,j) \in E} P_{ij}^{st} - \sum_{j:(j,i) \in E} P_{ji}^{st} = \begin{cases} 1 & i = s \\ -1 & i = t \\ 0 & \text{otherwise} \end{cases}$$

$$\forall (s,t) \in E_L, c \in C, (i,j) \in E$$

$$R^{cd} \geq A_d^{sc} \quad \forall d \in D, c \in C, s \in V_L$$

$$\sum_{d \in D} R^{cd} \leq K \quad \forall c \in C$$

$$\sum_{(i,j) \in Z} P_{ij}^{st} / M \leq \alpha_z^{st} \leq \sum_{(i,j) \in Z} P_{ij}^{st} \quad \forall (s,t) \in E_L, z \in Z$$

$$\sum_{(s,t) \in CS(S, V_L - S)} P_{ij}^{st} + P_{ji}^{st} < |CS(S, V_L - S)|$$

$$\forall S \subset V_L, (i,j) \in E$$

$$\sum_{(s,t)} P_{ij}^{st} \leq W \quad \forall (i,j) \in E$$

$$F_{st}^{u,c} \leq \alpha_z^{st} \quad \forall u \in V_L, c \in C, (s,t) \in E_L, z \in E$$

$$T_{st}^{u,c} \leq F_{st}^{u,c} + F_{st}^{u,c} - 1 \quad \forall u \in V_L, c \in C, (s,t) \in E_L; z_i, z_j \in E$$

$$T_{st}^{u,c} \leq F_{st}^{u,c}, F_{st}^{u,c} \quad \forall u \in V_L, c \in C, (s,t) \in E_L; z_i, z_j \in E$$

$$\sum_{j:(i,j) \in E} T_{st}^{u,c} - \sum_{j:(j,i) \in E} T_{ts}^{u,c} \begin{cases} 1 - A_d^{sc} & u = s \cap u \in D \\ 1 & u = s \\ -A_d^{sc} & s \in D \\ 0 & \text{otherwise} \end{cases}$$

$$\forall (s,t) \in E_L, c \in C, (i,j) \in E, z_i, z_j \in E$$

Mapping virtual links
to physical links



INTEGER LINEAR PROGRAMMING FORMULATION (NC1+CC2)

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$$\sum_{(s,t) \in CS(S, V_L - S)} P_{ij}^{st} + P_{ji}^{st} < |CS(S, V_L - S)|$$

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$$\sum_{(s,t)} P_{ij}^{st} \leq W \quad \forall (i,j) \in E$$

$$\forall (s,t) \in E_L, c \in C, (i,j) \in E, z_i, z_j \in E$$

Content placement



INTEGER LINEAR PROGRAMMING FORMULATION (NC1+CC2)

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$$\forall (s,t) \in E_L, c \in C, (i,j) \in E, z_i, z_j \in E$$

Select disrupted virtual links



INTEGER LINEAR PROGRAMMING FORMULATION (NC1+CC2)

$$\sum_{j:(i,j) \in E} P_{ij}^{st} - \sum_{j:(j,i) \in E} P_{ji}^{st} = \begin{cases} 1 & i = s \\ -1 & i = t \\ 0 & \text{otherwise} \end{cases}$$

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$$T_{st}^{u,c} z_i z_j \leq F_{st}^{u,c} z_i + F_{st}^{u,c} z_j - 1 \quad \forall u \in V_L, c \in C, (s,t) \in E_L; z_i, z_j \in E$$

$$R^{cd} \geq A_d^{sc} \quad \forall d \in D, c \in C, s \in V_L$$

$$T_{st}^{u,c} z_i z_j \leq F_{st}^{u,c} z_i, F_{st}^{u,c} z_j \quad \forall u \in V_L, c \in C, (s,t) \in E_L; z_i, z_j \in E$$

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$$\sum_{(i,j) \in Z} P_{ij}^{st} / M \leq \alpha_z^{st} \leq \sum_{(i,j) \in Z} P_{ij}^{st} \quad \forall (s,t) \in E_L, z \in Z$$

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$$\sum_{(s,t) \in CS(S, V_L - S)} P_{ij}^{st} + P_{ji}^{st} < |CS(S, V_L - S)|$$

$$\forall S \subset V_L, (i,j) \in E$$

$$\sum_{(s,t)} P_{ij}^{st} \leq W \quad \forall (i,j) \in E$$

Survivable mapping constraint



INTEGER LINEAR PROGRAMMING FORMULATION (NC1+CC2)

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$$T_{st}^{u,c} z_i z_j \leq F_{st}^{u,c} z_i + F_{st}^{u,c} z_j - 1 \quad \forall u \in V_L, c \in C, (s,t) \in E_L; z_i, z_j \in E$$

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$$\forall S \subset V_L, (i,j) \in E$$

$$\sum_{(s,t)} P_{ij}^{st} \leq W \quad \forall (i,j) \in E$$

Capacity constraint



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Reachability constraints



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$$R^{cd} \geq A_d^{sc} \quad \forall d \in D, c \in C, s \in V_L$$

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$$\sum_{d \in D} R^{cd} \leq K \quad \forall c \in C$$

$$\sum_{j:(i,j) \in E} T_{st}^{u,c} - \sum_{j:(j,i) \in E} T_{ts}^{u,c} \begin{cases} 1 - A_d^{sc} & u = s \cap u \in D \\ 1 & u = s \\ -A_d^{sc} & s \in D \\ 0 & \text{otherwise} \end{cases}$$

$$\sum_{(i,j) \in Z} P_{ij}^{st} / M \leq \alpha_z^{st} \leq \sum_{(i,j) \in Z} P_{ij}^{st} \quad \forall (s,t) \in E_L, z \in Z$$

$$\forall (s,t) \in E_L, c \in C, (i,j) \in E, z_i, z_j \in E$$

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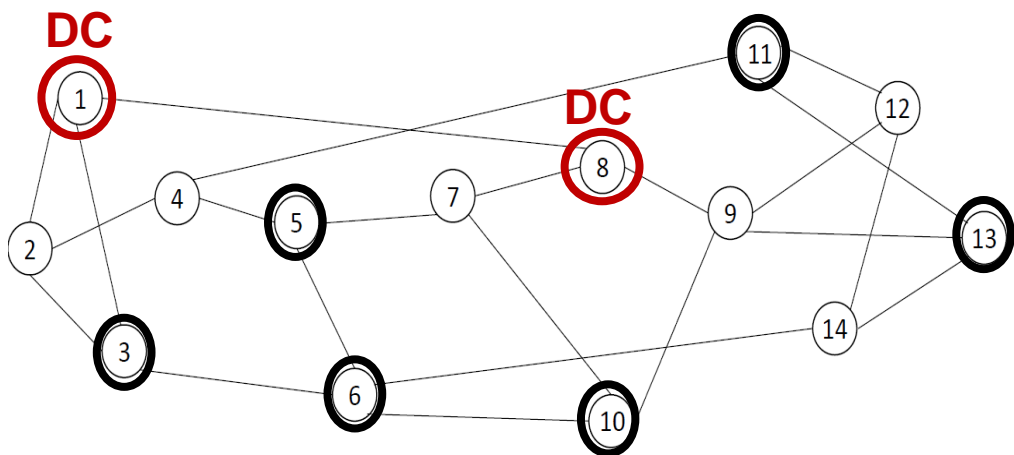
Objective function:

Minimize $\sum_{(s,t) \in E_L} \sum_{(i,j) \in E} P_{ij}^{st}$

$$\sum_{(s,t)} P_{ij}^{st} \leq W \quad \forall (i,j) \in E$$

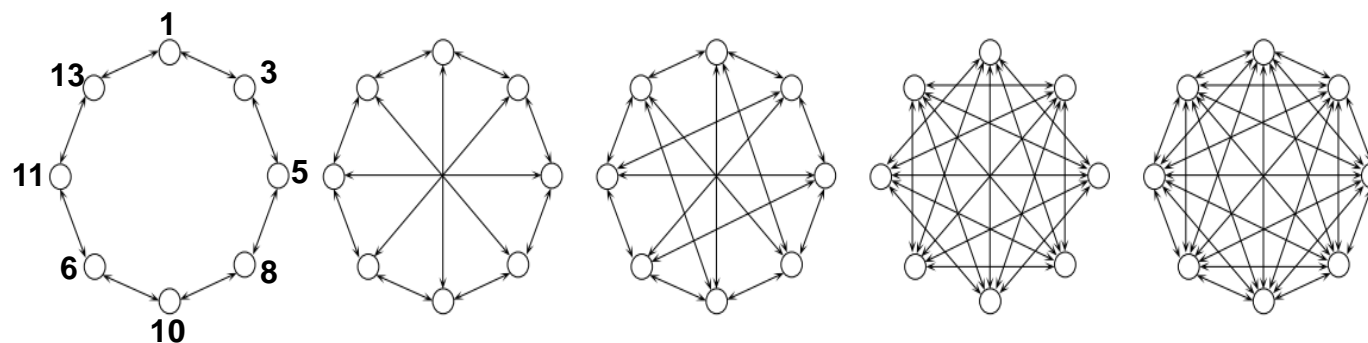


Physical topology



NSFNET (14 Nodes, 22 bidirectional links)

Logical topologies



0.29

0.47

0.57

0.71

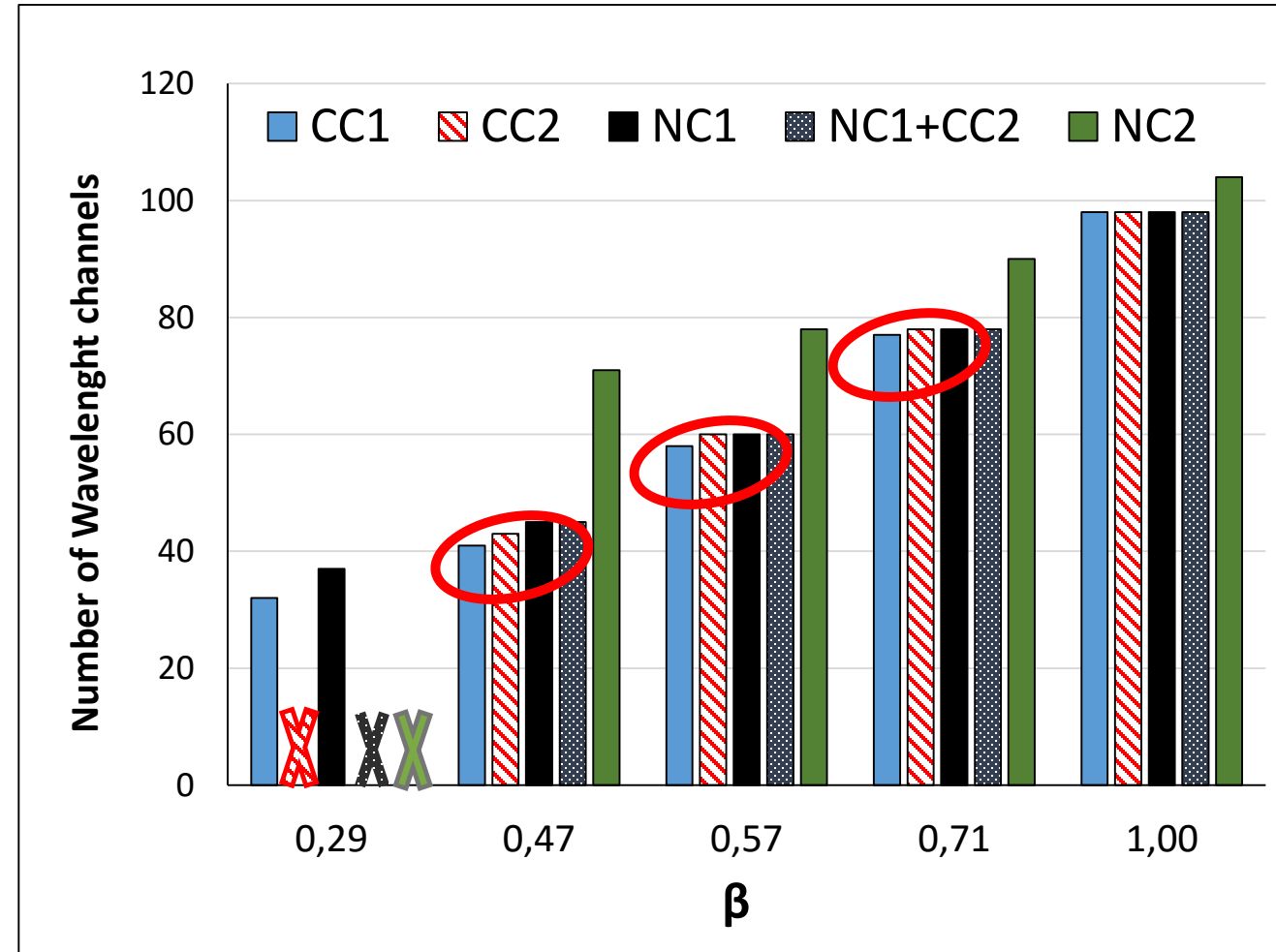
1

Connectivity degree (β)



RESOURCE USAGE COMPARISON

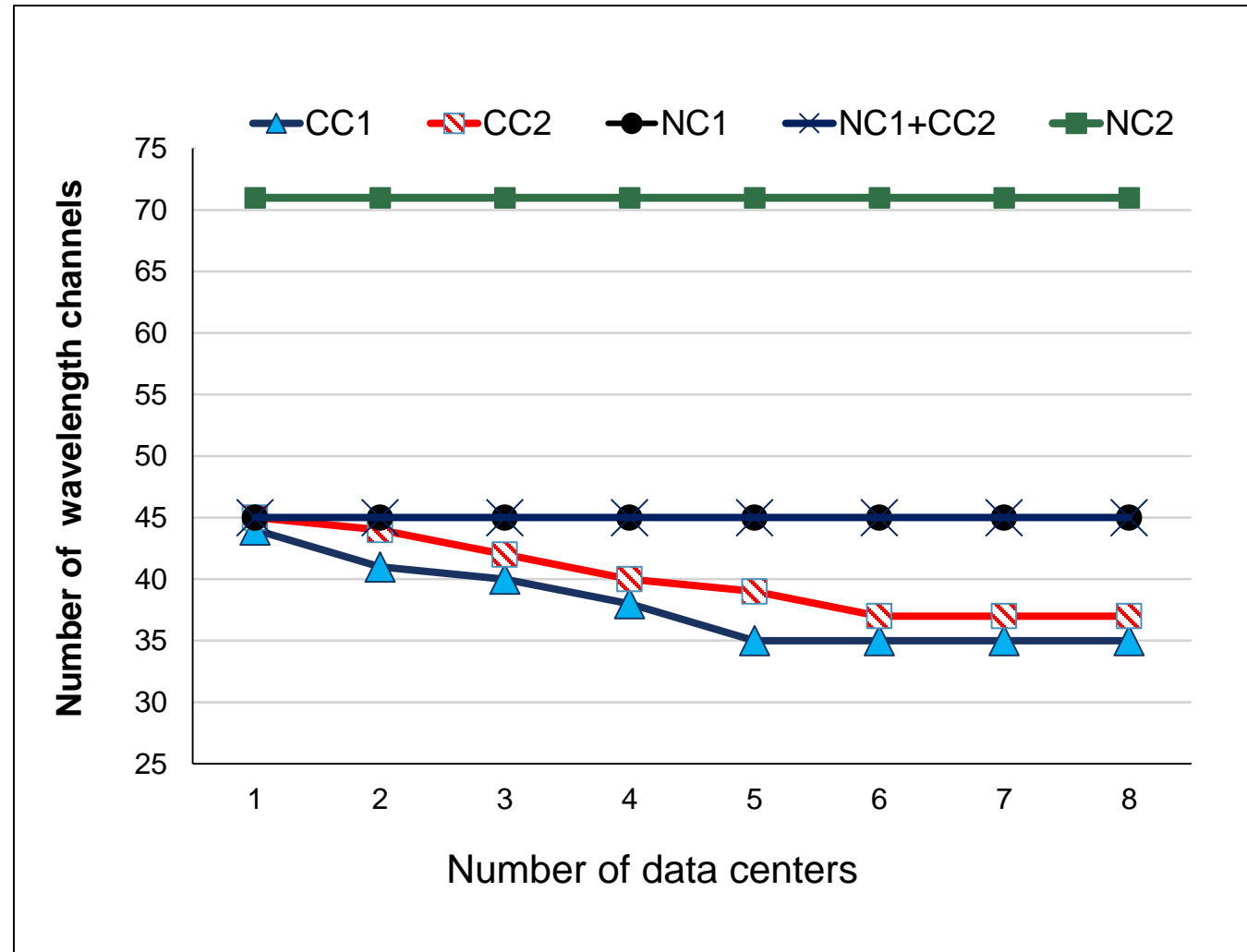
- ❑ Disaster locations: **All Physical links**
- ❑ Limited set of potential data center locations
- ❑ Number of data centers: **2**
- ❑ Limited number of wavelength per link





Assumptions:

- Logical topology $\beta = 0.47$
- All nodes are assumed to hold a data center





CONCLUSIONS

- For the proposed logical topologies, we show that it is possible to guarantee content connectivity with minimum resources and with a limited number of data centers.

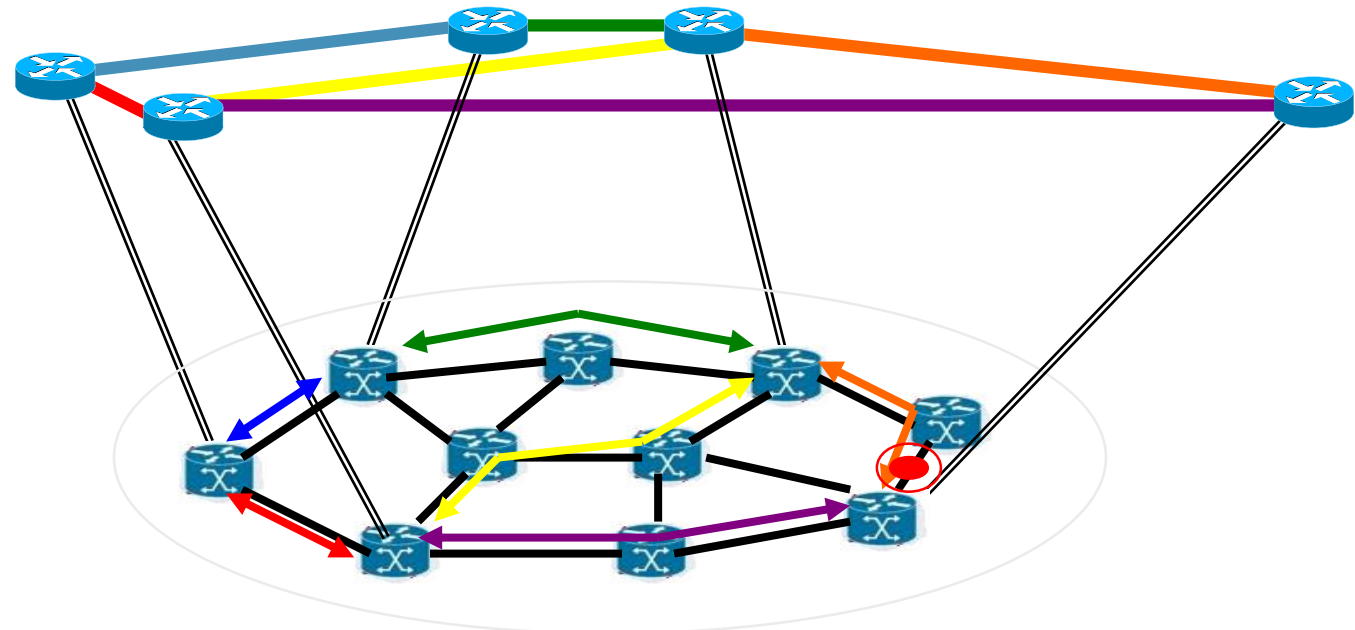
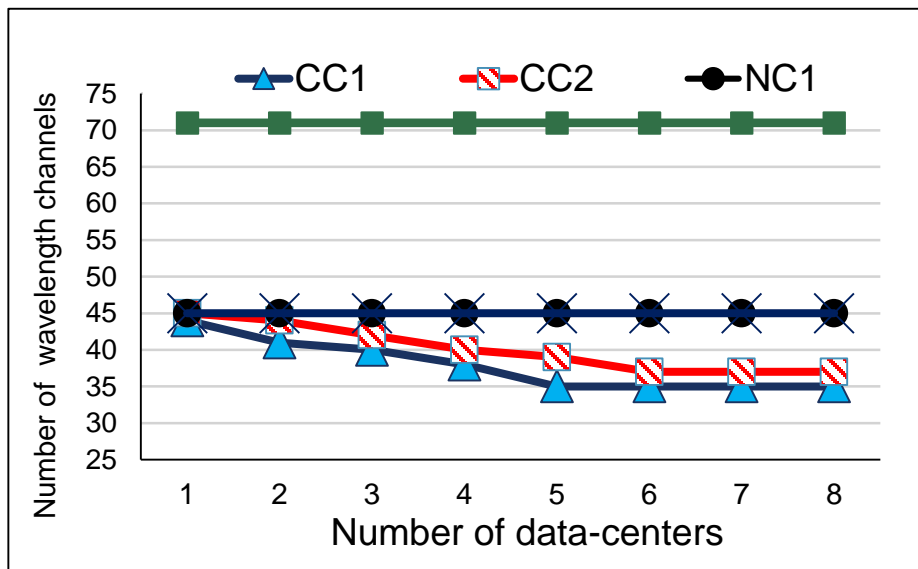
- CC1 requires almost the same amount of resources as NC1 for all design scenarios
 - No impact of the logical topology

- It is possible to guarantee continuity of the service after double-link failures with less resources
 - NC1 + CC2 requires 53% less resources with respect to NC2 on 8 node mesh virtual network with average node degree equal to 3
 - Especially when NC2 is not applicable



FUTURE WORKS

- Investigate theoretical bounds of content connectivity
- Design an heuristic to solve the problem for large instances
- Solve the problem under new constraints such as latency, storage capacity and content categorization





Thank you